



School of Business,  
Economics and Law  
GÖTEBORG UNIVERSITY

# Auctions

Johan Stennek

# Auctions

- Examples
  - Antiques, fine arts
  - Houses, apartments, land
  - Government bonds, bankrupt assets
  - Government contracts (roads)
  - Radio frequencies

# Auctions

## Why use auction?

- Seller's goal
  - Maximize revenues (you selling your apartment)
  - Efficient use (Government selling radio spectrum)
- Problem
  - Seller doesn't know what people are willing to pay
    - What is the highest valuation?
    - Who has it?
- Solution
  - Buyer *claiming* highest valuation gets the good
  - And will pay accordingly
- Auction = Mechanism to extract information

# Auctions

But are auctions a good solution?

- **Efficiency**
  - IF: people really “tell the truth” = bid their valuations
  - THEN: good will be allocated correctly
- **Revenues**
  - IF: people really “tell the truth” = bid their valuations
  - THEN: price will be high (efficiency & extract WTP)
- **Question: Do people “tell the truth”?**
  - Need to study bidding behavior

# Auctions

- Bidding behavior turns out to depend on:
  - Exact rules of the auction (Auction design)
  - How buyer's valuations are related (Type of uncertainty)

# Auctions

## 4 Designs

- Sealed bid, second price (“Vickrey”)
    - Simultaneous
    - Winner pays second bid
  - Sealed bid, first price
    - Simultaneous
    - Winner pays own bid
  - English (“open cry”)
    - Sequential + perf. info
    - Ascending bids
  - Dutch
    - Sequential + perf. info
    - Descending offers
- We will only study most common*

# Auctions

## Types of Uncertainty

- Private value
  - Different buyers have different values
- Common value
  - Same value to all buyers
  - But different buyers have different information

**We will only study private value**

# English Auction



# English Auction

- Assume
  - One indivisible unit of the good
  - Two bidders
- Information
  - Bidders get to know own valuations,  $v_1$  and  $v_2$
  - Then the bidding game starts
- Bidding rules: a simple model
  - Players take turns bidding
  - Whenever one player does not bid at least €1 more, the good is sold to the current bid

# English Auction

- Outcome
  - Winner = Highest bidder
  - Price = Highest bid

# Second-Price Sealed-Bids

- Utility

$$u_i = \begin{cases} v_i - b_j & \text{if winning} \\ 0 & \text{otherwise} \end{cases}$$

# English Auction

- Define: “marginal increases strategy” for  $i$ 
  - If current bid  $<$  valuation, raise by €1
  - If current bid  $>$  valuation, stop bidding
- Formally
  - IF:  $b_{jt-1} + 1 \leq v_i$ , THEN: bid  $b_{it} = b_{jt-1} + 1$
  - IF:  $b_{jt-1} + 1 > v_i$ , THEN: stop bidding
- Claim
  - This strategy is optimal (actually, dominant)

# English Auction

- Sketch of proof
  - If  $b_{2t-1} < v_1$ 
    - **Outbid:** Positive utility with (weakly) positive probability
    - **Withdraw:**  $u_1 = 0$  for sure
    - No reason to raise by more than €1
  - If  $b_{2t-1} \geq v_1$ 
    - **Withdraw:**  $u_1 = 0$  for sure
    - **Outbid:** Negative utility with (weakly) positive probability
  - Note - dominance
    - Above strategy optimal
    - no matter how  $b_{2t-1}$  selected

# English Auction

- Outcome
  - Q: “Truth telling”?
    - Sort of...
    - people keep raising the price until the bid is equal to their valuation (or nobody else continues to bid)
  - Q: Who gets the good?

# English Auction

- Outcome
  - “Truth telling”
  - Efficiency
    - Bidder with highest valuation wins the good
  - Q: Who gets the surplus?

# English Auction

- Outcome
  - “Truth telling”
  - Efficiency
    - Bidder with highest valuation wins the good
  - Surplus-sharing
    - $p = SHV$  (sometimes  $p = SHV + 1$ )



# First-Price Sealed-Bids Auction

# First-Price Sealed-Bids

- Rules
  - Simultaneous bids (= sealed bids)
  - Winner pays his bid (= first price)

# First-Price Sealed-Bids

- Trade-off
  - Higher bid  $\rightarrow$  Higher probability of winning
  - Higher bid  $\rightarrow$  Higher price

# First-Price Sealed-Bids

- Simplification

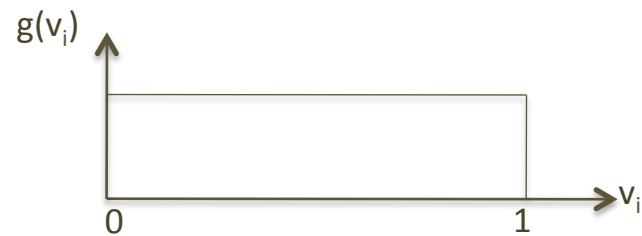
- Two bidders:  $v_1, v_2$

- $v_i$  uniformly distributed over  $[0, 1]$

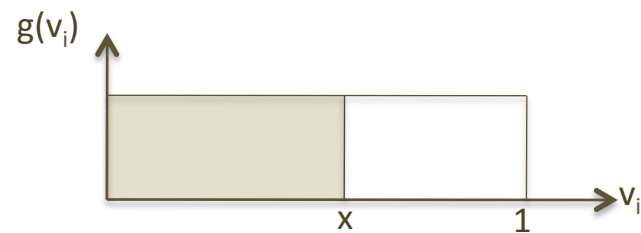


# First-Price Sealed-Bids

- Q: Probability that  $v_i < x$ ?



- A:  $\text{Prob}(v_i < x) = x$



# First-Price Sealed-Bids

- Payoff = expected utility

- $E\pi_1(b_1) = (v_1 - b_1) \Pr(\text{win}) + 0 \Pr(\text{lose})$

- $E\pi_1(b_1) = (v_1 - b_1) \Pr(b_1 > b_2)$

**Depends on**

$b_1$  = own choice

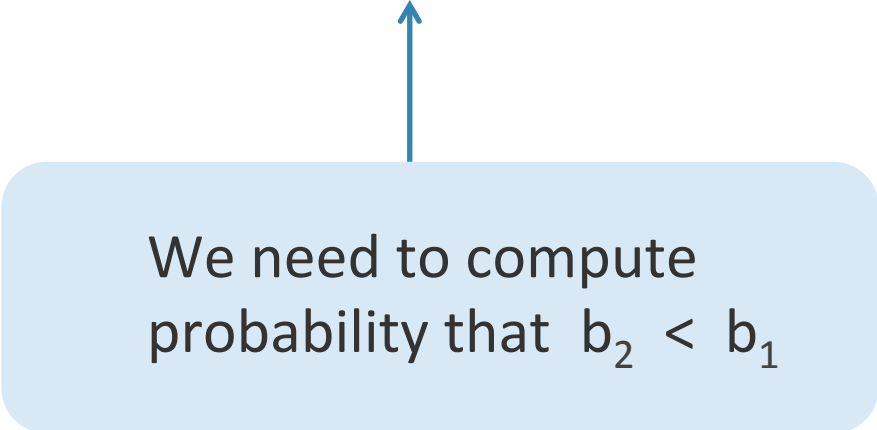
$b_2$  = random variable

# First-Price Sealed-Bids

- Payoff = expected utility

- $E\pi_1(b_1) = (v_1 - b_1) \Pr(\text{win}) + 0 \Pr(\text{lose})$

- $E\pi_1(b_1) = (v_1 - b_1) \Pr(b_1 > b_2)$



We need to compute  
probability that  $b_2 < b_1$

# First-Price Sealed-Bids

- Payoff = expected utility

- $E\pi_1(b_1) = (v_1 - b_1) \Pr(\text{win}) + 0 \Pr(\text{lose})$

- $E\pi_1(b_1) = (v_1 - b_1) \Pr(b_1 > b_2)$

Simplifying assumption:

$$b_2 = z \cdot v_2$$



# First-Price Sealed-Bids

- Payoff = expected utility

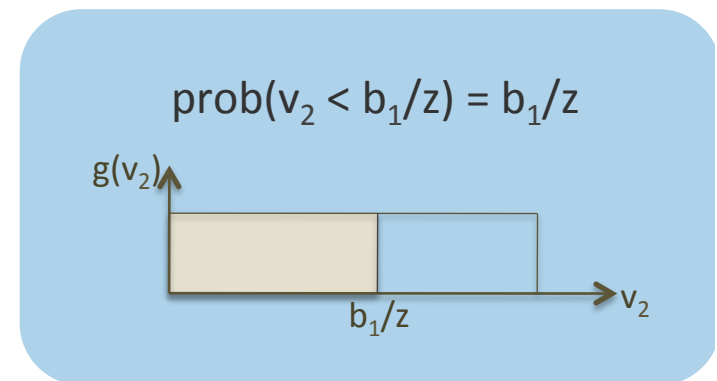
- $E\pi_1(b_1) = (v_1 - b_1) \Pr(\text{win}) + 0 \Pr(\text{lose})$

- $E\pi_1(b_1) = (v_1 - b_1) \Pr(b_1 > b_2)$

- $E\pi_1(b_1) = (v_1 - b_1) \Pr(b_1 > z \cdot v_2)$

- $E\pi_1(b_1) = (v_1 - b_1) \Pr(v_2 < b_1/z)$

- $E\pi_1(b_1) = (v_1 - b_1) (b_1/z)$



# First-Price Sealed-Bids

- Conclusion
  - IF:  $b_2 = z \cdot v_2$
  - THEN:  $E\pi_1(b_1) = (v_1 - b_1) (b_1/z)$
- Q: What is player 1's best reply?

# First-Price Sealed-Bids

- What is 1's best reply?

- $E\pi_1(b_1) = (v_1 - b_1) (b_1/z)$

- FOC:  $(-1) (b_1/z) + (v_1 - b_1) (1/z) = 0$

Utility if winning \* Increased probability of winning



# First-Price Sealed-Bids

- Assume

- $B_2(v_2) = z v_2$

- What is 1's best reply?

- $E\pi_1(b_1) = (v_1 - b_1) (b_1/z)$

- FOC:  $(-1) (b_1/z) + (v_1 - b_1) (1/z) = 0$



Decreased utility \* probability of winning

# First-Price Sealed-Bids

## Proof

- Assume
  - $B_2(v_2) = z v_2$
- What is 1's best reply?
  - $E\pi_1(b_1) = (v_1 - b_1) (b_1/z)$
  - FOC:  $-(b_1/z) + (v_1 - b_1)/z = 0$
  - Solve:  $b_1 = \frac{1}{2} \cdot v_1$

# First-Price Sealed-Bids

## Proof

- Conclusion

- IF: Bidder 2 uses a linear strategy:  $B_2(v_2) = z \cdot v_2$

- THEN: Best reply for bidder 1:  $B_1(v_1) = \frac{1}{2} \cdot v_1$

- Note

- Since  $\frac{1}{2} \cdot v_1$  is linear

- Since players are symmetric

- Both bidding  $b_i = \frac{1}{2} \cdot v_i$  is a Nash equilibrium of a game where the strategy for each player is to choose some function  $B_i(v_i)$ .

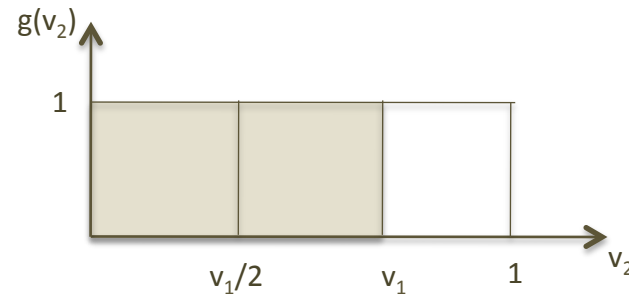
# First-Price Sealed-Bids

- Interpretation
  - Why bid  $\frac{1}{2} v$  ?
- Answer 1
  - Optimal balance between
    - probability of winning
    - price in case of winning

# First-Price Sealed-Bids

- Interpretation

- But why exactly  $\frac{1}{2}$  ?



- Answer 2

- Assume you have highest valuation

- Q: What is the expected second highest valuation?

- Winner bids expected wtp of competitor  
=> competitor no incentive to bid more



# First-Price Sealed-Bids

- Remark
  - With more bidders, expected second highest wtp is closer to highest wtp
  - Bid larger share of wtp
  - As  $n \rightarrow \infty$   $b \rightarrow wtp$

# First-Price Sealed-Bids

- Outcome
  - Q: Who gets the good?

# First-Price Sealed-Bids

- Outcome
  - Efficiency
    - Bidder with highest valuation wins the good
  - Q: Who gets the surplus?

# First-Price Sealed-Bids

- Outcome
  - Efficiency
    - Bidder with highest valuation wins the good
  - Surplus-sharing
    - $p = \frac{1}{2} HV$
  - Truth-telling?

# First-Price Sealed-Bids

- Outcome

- Efficiency

- Bidder with highest valuation wins the good

- Surplus-sharing

- $p = \frac{1}{2} HV$

- “Sort of truth-telling”

- Players actually reveal their valuation

# Game Theoretic “Details”

Auction = Game of Incomplete Information

# Game of Incomplete Information

- Game with incomplete information
  - Buyers don't know each others' valuations
    - Ada is not able to predict Ben's bid exactly
    - It depends on Ben's valuation of the object
  - How should Ada and Ben analyze the situation?

# Game of Incomplete Information

- Solution I: Change definition of strategy
  - Strategy = Function prescribing bid for every possible valuation a player may have
- Example of strategy
  - IF  $wtp = v_H$  THEN  $bid = b_H$
  - IF  $wtp = v_L$  THEN  $bid = b_L$
- Then, players able to
  - Predict rival's strategy, even if uncertainty about type and bid remains
  - Maximize expected payoff



# Game of Incomplete Information

- But why are strategies functions?
  - Ada knows she has high valuation,  $v_H$
  - Why should she choose strategy with instruction for  $v_L$ ?
- Answer
  - Ben doesn't know Ada's valuation. Could be  $v_H$  or  $v_L$
  - Ben must consider
    - What would Ada bid if  $v_H$
    - What would Ada bid if  $v_L$
  - To predict Ben's bid, Ada must also consider what she herself would have bid in case of  $v_L$

# Game of Incomplete Information

- Think of Ada's choice as two-step procedure
  1. Find optimal bid for all possible valuations:  
 $b^{\text{Ada}}(v_H)$  and  $b^{\text{Ada}}(v_L)$
  2. Select the relevant bid:  $b^{\text{Ada}}(v_H)$

# Game of Incomplete Information

- Solution II: Change definition of payoff
  - Payoff = expected utility

# Game of Incomplete Information

## Bayesian Nash Equilibrium

- Pair of strategies  $(b^{Ada}, b^{Ben})$  such that
- $b^{Ada}$  is a best reply to  $b^{Ben}$ 
  - $b^{Ada}(v_H)$  maximizes Ada's expected utility
    - If Ada's valuation is  $v_H$
    - Assuming Ben uses  $b^{Ben}$
  - $b^{Ada}(v_L)$  maximizes Ada's expected utility
    - If Ada's valuation is  $v_L$
    - Assuming Ben uses  $b^{Ben}$
- $b^{Ben}$  is a best reply to  $b^{Ada}$

Most fundamental result of  
auction theory

# Fundamental result

Note 1: No individual knows who has the highest valuation

Note 2: But if people play the auction game  
⇒ person with highest valuation walks away with the good

No individual (even a dictator)  
could have implemented the efficient allocation,  
since nobody has sufficient information

But the market mechanism actually solves the maximization problem

May say the market aggregates information

- \* must use all the information to solve the max-problem
- \* despite the fact that it is scattered

# Fundamental result

- Laboratory experiments
  - It works! (Vernon Smith)
  - Also double auctions
  - Even with “few” buyers and “few” sellers market quickly converges to competitive price
  - NB: must use laboratory to know people’s valuations

# Fundamental result

Sure, it is not perfect...

...there is also market failure...

- Coordination (mis-pricing; recessions)
- Double coincidence of wants (kidneys, apartments)
- Externalities (global warming; telecom)
- Public goods (R&D; legal system to enforce all contracts)
- Market power (medicines; district heating)
- Incomplete information (cars, insurance, labor, credit)

...and an uneven distribution of wealth



# Fundamental result

- But even public policies to correct market failure use markets to aggregate information
  - Cap and trade
  - Public procurement

# Comparison of Auction Designs (Revenues)

# Comparison of Designs

## Question 1

- Which auction gives the highest expected price?
  - FPSB (and Dutch):  $p = \frac{1}{2} HV$
  - English (and SPSB):  $p = SHV$       Recall:  $E(SHV) = \frac{1}{2} HV$

# Comparison of Designs

## Answer 1

- Expected Revenue Equivalence Thm.  
(Vickrey, 1961)
  - All four auctions give the same *expected* price

# Comparison of Designs

## Question 2

- Is there any other way to sell the goods which would give a higher expected profit?
  - Lots of different possible ways
    - Bargaining
    - Other auction formats
    - Strange games

# Comparison of Designs

## Answer 2

- Generalization of Revenue Equivalence Thm
  - No!
  - This is example of “mechanism design” and uses the “revelation principle” (Leonid Hurwicz, Eric Maskin, Roger Myerson)